

# Kotlin language specification

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## Chapter 9

# Operator overloading

Some syntax forms in Kotlin are *defined by convention*, meaning that their semantics are defined through syntactic expansion of one syntax form into another syntax form.

Particular cases of definition by convention include:

- Arithmetic and comparison operators;
- `invoke` convention;
- Operator-form [assignments](#);
- [For-loop statements](#);
- [Delegated properties](#);
- [Destructuring declarations](#).

Important: another case of definition by convention is [safe navigation](#), which is covered in more detail in its respective section.

There are several points shared among all the syntax forms defined using definition by convention:

- The expansions are hygienic: if they introduce new identifiers that were not present in original syntax, all such identifiers are not accessible outside the expansion and cannot clash with any other declarations in the program;
- The expressions captured by an expansion are using *call-by-need* evaluation strategy, meaning that they are evaluated only once during first usage specified in the expansion even if the expansion itself has more than one usage of such an expression;
- An expansion may lead to another expansion, following the same rules;
- All call expressions that are produced by expansion are only allowed to use operator functions.

This expansion of a particular syntax form to a different piece of code is usually defined in the terms of *operator* functions.

Operator functions are function which are [declared](#) with a special keyword `operator` and are not different from regular functions when called via [function calls](#). However, operator functions can also be used in definition by convention.

Note: it is not important whether an operator function is a member or an extension, nor whether it is suspending or not. The only requirements are the ones listed in the respected sections.

For example, for an operator form `a + b` where `a` is of type `A` and `b` is of type `B` any of the following function definitions are applicable:

```
class A {
  // member function
  operator fun plus(b: B) = ...
  // suspending member function
  suspend operator fun plus(b: B) = ...
}

// extension function
operator fun A.plus(b: B) = ...
// suspending extension function
suspend operator fun A.plus(b: B) = ...
```

Assuming additional implicit receiver of this type is available, it may also be an extension defined in another type:

```
object Ctx {
  // extension that is a member of some context type
  operator fun A.plus(b: B) = ...

  fun add(a: A, b: B) = a + b
}
```

Note: different platforms may add additional criteria on whether a function may be considered a suitable candidate for operator convention.

The details of individual expansions are available in the sections of their respective operators, here we would like to describe how they *interoperate*.

For example, take the following declarations:

```
class A {
  operator fun inc(): A { ... }
}

object B {
  operator fun get(i: Int): A { ... }
  operator fun set(i: Int, value: A) { ... }
}
```

```
object C {
  operator fun get(i: Int): B { ... }
}
```

The expression `C[0][0]++` is expanded (see the [Expressions](#) section for details) using the following rules:

- The operations are expanded in order of their priority.
- First, the [increment operator](#) is expanded, resulting in:

```
C[0][0] = C[0][0].inc()
```

- Second, the [assignment](#) to an indexing expression (produced by the previous expansion) is expanded, resulting in:

```
C[0].set(C[0][0].inc())
```

- Third, the [indexing expressions](#) are expanded, resulting in:

```
C.get(0).set(C.get(0).get(0).inc())
```

Important: although the resulting expression contains several instances of the subexpression `C.get(0)`, as all these instances were created from the same original syntax form, the subexpression is evaluated only once, making this code roughly equivalent to:

```
val $tmp = C.get(0)
$tmp.set($tmp.get(0).inc())
```

## 9.1 Destructuring declarations

A special case of definition by convention is the destructuring declaration of properties, which is available for [local properties](#), parameters of [lambda literals](#) and the iteration variable of [for-loops](#). See the corresponding sections for particular syntax.

This convention allows to introduce a number (one or more) of properties in the place of one by immediately *destructuring* the property during construction. The immediate value (that is, the initializing expression of the local property, the value acquired from the operator convention of a for-loop statement, or an argument passed into a lambda body) is assigned to a number of placeholders  $p_0, \dots, p_N$  where each placeholder is either an identifier or a special ignoring placeholder `_` (note that `_` is not a valid identifier in Kotlin). For each identifier the corresponding operator function `componentK` with  $K$  being equal to the position of the placeholder in the declaration (**starting from 1**) is called without arguments and the result is assigned to a fresh value referred to as the identifier used. For each ignoring placeholder, no calls are performed and nothing is assigned. Each placeholder may be provided with an optional type signature

$T_M$  which is used in [type inference](#) as any property type would. Note that an ignoring placeholder may also be provided with a type signature, in which case although the call to corresponding `componentM` function is not performed, it still must be checked for function applicability during type inference.

Examples:

```
val (x: A, _, z) = f()
```

is expanded to

```
val $tmp = f()
val x: A = $tmp.component1()
val z = $tmp.component3()
```

where `component1` and `component3` are suitable operator functions available on the value returned by `f()`

```
for((x: A, _, z) in f()) { ... }
```

is expanded to (as per for-loop expansion)

```
when(val $iterator = f().iterator()) {
  else -> while ($iterator.hasNext()) {
    val $tmp = $iterator.next()
    val x: A = $tmp.component1()
    val z = $tmp.component3()
    ...
  }
}
```

where `iterator()`, `next()`, `hasNext()`, `component1()` and `component3` are all suitable operator functions available on their respective receivers.

```
foo { (x: A, _, z) -> ... }
```

is expanded to

```
foo { $tmp ->
  val x: A = $tmp.component1()
  val z = $tmp.component3()
  ...
}
```

where `component1()` and `component3` are all suitable operator functions available on the value of lambda argument.